

Dual Dynamic Programming for Long-Term Receding-Horizon Hydro Scheduling with Head Effects

<u>Benjamin Flamm</u>, Annika Eichler, Joseph Warrington, John Lygeros Automatic Control Laboratory, ETH Zurich

SINTEF Hydro Power Scheduling Workshop 2018



Motivation: Long-term control of energy storage

- Nonlinear, complex models
- Receding horizon setting
- Applications:
 - Seasonal demand
 - Peak shaving
 - Price arbitrage







Seasonal storage problem

• Want to solve long-term nonlinear problem of the form (NLP)

$$\min_{x,u} \sum_{t=0}^{T-1} g_t(x_t, u_t) + g_T(x_T)$$

s.t. $x_{t+1} = f_t(x_t, u_t),$ $t = 0, \dots, T-1$
 $(u_t, x_t) \in \mathcal{Z}_t,$ $t = 0, \dots, T-1$
 x_0 given.

... but this can be computationally intractable.



Split-horizon approximation

• To make problem tractable, approximate second part of horizon: (NLP+LP)

$$\min_{x,u} \sum_{t=0}^{T_1-1} g_t(x_t, u_t) + \sum_{t=T_1}^{T-1} \left(c_t^{\top} x_t + d_t^{\top} u_t \right) + c_T^{\top} x_T$$

s.t. $x_{t+1} = \begin{cases} f_t(x_t, u_t), & t = 0, \dots, T_1 - 1 \\ A_t x_t + B_t u_t, & t = T_1, \dots, T - 1 \end{cases}$
 $(u_t, x_t) \in \mathcal{Z}_t, & t = 0, \dots, T_1 - 1 \\ E_t x_t + F_t u_t \le h_t, & t = T_1, \dots, T - 1 \end{cases}$
 x_0 given.



Split-horizon approximation, continued

Split problem into nonlinear part (NLP1):

$$\min_{\substack{u_0,\ldots,u_{T_1-1}\\x_1,\ldots,x_{T_1}}} \sum_{t=0}^{T_1-1} g_t(x_t, u_t) + \tilde{G}_{T_1}(x_{T_1})$$

s.t. $x_{t+1} = f_t(x_t, u_t), \quad t = 0, \ldots, T_1 - 1$
 $(u_t, x_t) \in \mathcal{Z}_t, \qquad t = 0, \ldots, T_1 - 1$
 x_0 given.

• And linear part (LP2): $\tilde{G}_{T_{1}}(x_{T_{1}}) = \min_{\substack{u_{T_{1}},...,u_{T-1}\\x_{T_{1}+1},...,x_{T}}} \sum_{t=T_{1}}^{T-1} \left(c_{t}^{\top}x_{t} + d_{t}^{\top}u_{t}\right) + c_{T}^{\top}x_{T}$ s.t. $x_{t+1} = A_{t}x_{t} + B_{t}u_{t},$ $E_{t}x_{t} + F_{t}u_{t} \leq h_{t},$ $t = T_{1}, \ldots, T-1.$



Complete recourse assumption

Assumption 1: The linear second-stage problem (LP2) is feasible for all x_{T_1} in the domain of the first-stage problem (NLP1).



Benders decomposition review

• Want to find value function for (LP2), as function of terminal state x_{T_1}



• Approximate value function: $\tilde{G}_{T_1}^{(j)}(x_{T_1}) = \max_{i=1,\dots,j} \{a_i^\top x_{T_1} + b_i\}$



Pereira and Pinto, 1991

Benders decomposition bounds

Given feasible solutions for each stage, can form

$$LB = \sum_{t=0}^{T_1 - 1} g_t(x_t, u_t) + \tilde{G}_{T_1}^{(j)}(x_{T_1})$$
$$UB = \sum_{t=0}^{T_1 - 1} g_t(x_t, u_t) + \sum_{t=T_1}^{T - 1} \left(c_t^\top x_t + d_t^\top u_t \right) + c_T^\top x_T$$



Dual dynamic programming algorithm with approximate solutions of nonlinear first stage

Nonlinear first stage (NLP1):

$$\min_{\substack{u_0, \dots, u_{T_1-1} \\ x_1, \dots, x_{T_1}}} \sum_{t=0}^{T_1-1} g_t(x_t, u_t) + G_{T_1}(x_{T_1})$$

s.t. $x_{t+1} = f_t(x_t, u_t), \quad t = 0, \dots, T_1 - 1$
 $(u_t, x_t) \in \mathcal{Z}_t, \qquad t = 0, \dots, T_1 - 1$
 x_0 given.

• Linear second stage (LP2): $G_{T_1}(x_{T_1}) = \min_{\substack{u_{T_1}, \dots, u_{T-1} \\ x_{T_1+1}, \dots, x_T}} \sum_{t=T_1}^{T-1} \left(c_t^\top x_t + d_t^\top u_t \right) + c_T^\top x_T$ s.t. $x_{t+1} = A_t x_t + B_t u_t$, $E_t x_t + F_t u_t \le h_t$, $t = T_1, \dots, T-1$.

Algorithm 1 Split-horizon DDP with nonlinear first horizon

while
$$(UB - LB > \epsilon)$$
 do

1. Solve (NLP1) using
$$\tilde{G}_{T_1}^{(j)}(x_{T_1}) = \max_{i=1,...,j} \{a_i^\top x_{T_1} + b_i\};$$

- 2. Generate LB using x, u from step 1 and current $\tilde{G}_{T_1}^{(j)}(x_{T_1})$;
- 3. Solve (LP2) using x_{T_1} from step 1;
- 4. Generate UB using x, u from steps 1 and 3;

5. Improve estimate of $G_{T_1}(x_{T_1})$ using dual variables from step 3; 6. j = j + 1; end

Benjamin Flamm | Sep. 12, 2018 | 9

Dual dynamic programming algorithm convergence

Theorem 1: When using Algorithm 1 to solve a split-horizon problem of the form (NLP+LP), the algorithm will terminate in a **finite number of iterations**, returning a **feasible** solution of the split-horizon problem. In addition, two special cases apply:

- (a) If the nonconvex first stage is solved in step 1 to **local optimality**, then Algorithm 1 will return a **locally-optimal** solution to the two-stage problem.
- (b) If the nonconvex first stage is solved in step 1 to **global optimality**, then Algorithm 1 solves the two-stage problem to **global optimality**.

Proof sketch: Since second stage is linear, finite number of lower-bounding hyperplanes. Eventually find all relevant hyperplanes to set LB = UB at given point.

$$LB = \sum_{t=0}^{T_1 - 1} g_t(x_t, u_t) + \tilde{G}_{T_1}^{(j)}(x_{T_1})$$
$$UB = \sum_{t=0}^{T_1 - 1} g_t(x_t, u_t) + G(x_{T_1})$$



Benjamin Flamm | Sep. 12, 2018 | 10

Error bound on two-stage approximation

We restrict ourselves to the following tractable problem class:

Assumption 2: There exists a map M_t which takes a solution for (LP2) and produces a unique solution for the second stage of (NLP) such that

- (a) At each timestep, the objective of the LP is an underestimate of (NLP), with a maximum error of δ_t .
- (b) Through M_t , each feasible solution of (NLP) is mapped to by a feasible solution of (LP2).

In other words, any feasible second-stage solution of (NLP) is underapproximated by a feasible solution of (LP2), with a maximum underestimate of δ_t .



Theorems on suboptimality

• **Theorem 2**: When using Algorithm 1 to solve (NLP+LP), if Assumption 2 holds, and step 1 is solved to ϵ -suboptimality each iteration, then the resulting solution $\underline{G}_{0}^{(J)}(x_{0})$ is such that:

$$\underline{G}_{0}^{(J)}(x_{0}) \leq G_{0}(x_{0}) \leq \underline{G}_{0}^{(J)}(x_{0}) + \epsilon + \sum_{t=T_{1}}^{T-1} \delta_{t}.$$

• **Theorem 3**: When using Algorithm 1 to solve (NLP+LP), if Assumption 2 holds, and step 1 is solved to ϵ -suboptimality each iteration, then when the resulting arguments $(\tilde{x}^{(J)}, \tilde{u}^{(J)})$ are evaluated in the objective of (NLP):

$$G_0(x_0) \leq \sum_{t=0}^{T_1-1} g_t(\tilde{x}_t^{(J)}, \tilde{u}_t^{(J)}) + G_{T_1}(\tilde{x}_{T_1}^{(J)}) \leq G_0(x_0) + \epsilon + \sum_{t=T_1}^{T-1} \delta_t.$$



Inspired by Guigues, 2018

Application: hydro reservoir scheduling with head effects

- System of multiple, interconnected reservoirs, with head-dependent power
- Objective: meeting electricity demand, with deterministic electricity spot prices





Application: sample hydro reservoir system parameters

- Five equally sized reservoirs
- 35 MW turbines and pumps with system topology as below
- Drainage time roughly 10 days (from full to empty at full power)
- No spillage constant amount of water in system
- Reservoirs are rectangular in profile
 - Head can vary by a factor of 5 (40 m to 200 m)



Application: hydro reservoir scheduling with head effects

Energy per unit volume depends on net head between reservoirs:

 $E^{i \to j}[t] = \alpha^{i \to j} + \beta^{i \to j} \left(\ell_i[t] - \ell_j[t] \right), \quad P^{i \to j}[t] = V^{i \to j}[t] E^{i \to j}[t]$

First stage exact NLP:

$$\begin{split} \min_{V} & \sum_{t=0}^{T-1} p[t] \Big(D[t] + \sum_{i=1}^{N} \sum_{j \in \mathcal{N}^{i \to j}} P^{i \to j}[t] \Big) + G(\ell[T]) \\ \text{s.t.} & P^{i \to j}[t] = V^{i \to j}[t] \left(\alpha^{i \to j} + \beta^{i \to j} \left(\ell_i[t] - \ell_j[t] \right) \right) \quad \forall \ j \in \mathcal{N}^{i \to j} \\ & \ell_i[t+1] = \ell_i[t] + \frac{1}{\gamma_i} \Big(\sum_{k \in \mathcal{N}^{\to i}} V^{k \to i}[t] - \sum_{j \in \mathcal{N}^{i \to j}} V^{i \to j}[t] \Big) \\ & \underline{\ell}_i[t] \leq \ell_i[t] \leq \bar{\ell}_i[t], \quad \ell_i[0] \text{ given} \\ & \underline{V}^{i \to j}[t] \leq V^{i \to j}[t] \leq \bar{V}^{i \to j} \quad \forall \ j \in \mathcal{N}^{i \to}[t] \\ & i = 1, \dots, N; \quad t = 0, \dots, T-1. \end{split}$$

Bilinear approximation: McCormick envelope

We replace the problematic constraint $P^{i \to j}[t] = V^{i \to j}[t] \left(\alpha^{i \to j} + \beta^{i \to j} \left(\ell_i[t] - \ell_j[t] \right) \right)$ with

$$P^{i \to j}[t] = V^{i \to j}[t] \ \alpha^{i \to j} + \beta^{i \to j}(\chi_i^{i \to j}[t] - \chi_j^{i \to j}[t])$$

where:

$$\begin{split} \chi_{i}^{i \to j}[t] &\geq V^{i \to j}[t]\bar{\ell}_{i} + \bar{V}^{i \to j}\ell_{i}[t] - \bar{V}^{i \to j}\bar{\ell}_{i}, \\ \chi_{i}^{i \to j}[t] &\geq V^{i \to j}[t]\underline{\ell}_{i} + \underline{V}^{i \to j}\ell_{i}[t] - \underline{V}^{i \to j}\underline{\ell}_{i}, \\ \chi_{i}^{i \to j}[t] &\leq V^{i \to j}[t]\bar{\ell}_{i} + \underline{V}^{i \to j}\ell_{i}[t] - \underline{V}^{i \to j}\bar{\ell}_{i}, \\ \chi_{i}^{i \to j}[t] &\leq V^{i \to j}[t]\underline{\ell}_{i} + \bar{V}^{i \to j}\ell_{i}[t] - \bar{V}^{i \to j}\underline{\ell}_{i}. \end{split}$$



Bilinear approximation: McCormick envelope







Improved bounds on McCormick envelopes

Set linearization bounds to worst-case combination of inflows and outflows





Successive linearization with McCormick envelopes

Set linearization bounds to envelope about previously-found trajectory



Successive linearization with McCormick envelopes

Objective improves rapidly with heuristic algorithm



Benjamin Flamm | Sep. 12, 2018 | 20

Experimental results: objective

- Compute objective over receding horizon for 20 days
 - 1 day split-horizon method is 4% better than constant efficiency approximation, but only 90% of objective found by successive linearization





Discussion

- For larger systems and longer time horizons, worst-case bounds are quite poor
 - Modeling exact near term is beneficial
- Solving exact problem scales exponentially in problem size
 - But longer linear stage is relatively cheap



Conclusion

- Solved long-term nonlinear problem via multistage approximation
- Extended DDP to case where first stage is nonlinear
- Provided bounds on suboptimality

Different nonlinear problem classes?





Thanks!

Reference:

B. Flamm, A. Eichler, J. Warrington, J. Lygeros, "Dual Dynamic Programming for Nonlinear Control Problems over Long Horizons," European Control Conference, 2018. <u>https://polybox.ethz.ch/index.php/s/woxUYLtSEKRIfKB</u>

